**R-Type Datapath**

x3 x2 x1 add

x2 Read register 1

x3 Read register 2

x1 **Registers**

Write

ALUSrc

register

Write data

RegWrite

Read data 1

Read data 2

**add**

4

ALU

Zero

**ALU** ALU

**M u x**

0

1

operation

Address

result

Write data

MemWrite

Read 1 data **M**

MemtoReg

**u x**

0

**Data memory**

32 **Imm** 64

Instruction

**Gen**

MemRead

32 x2 x1 load

x2 Read register 1

Read register 2

x1 **Registers**

Write

register

Write data

RegWrite

32

Read data 1

Read data 2

**add**

4

ALU

Zero

**ALU** ALU

result

**M u x**

ALUSrc

0

1

64

operation

Address

Write data

MemWrite

Read

1

MemtoReg

data **M u x**

0

**Data memory**

32 **Imm**

Instruction

**Gen**

MemRead

**Andi x2, x1, 25**

25 x1 x2 andi

Inst

Instruction

Word

x1 Read register 1

Read register 2

x2 **Registers**

Write

register

Write data

RegWrite

Read data 1

Read

data 2 **M u x**

ALUSrc

0

1

64

25

**and**

4

ALU

Zero

**ALU** ALU

result

operation

Address

Write data

MemWrite

Read 1 data **M**

MemtoReg

**u x**

0

**Data memory**

32 **Imm**

**Gen**

MemRead

32 x1 x2 store

x2 Read register 1

Read register 2

Read data 1

**add**

4

ALU

Zero

**ALU** ALU

result

operation

MemWrite

MemtoReg

Write

**Registers**

Read

ALUSrc

0

Address

Read 1

data

register

x1

Instruction

Write data

RegWrite

IW

32

data 2 **M u x**

1

64

32

**Imm**

**Gen**

Write data

**M u x**

0

**Data memory**

MemRead

offset x2 x1 beq

Branch target address is PC + Imm \* 2 (Slide #20)

PCSrc

0

4

Add

Read address

**PC**

x2 Read register 1

x1 Read register 2

Read data 1

**Shift**

**left 1**

ALUSrc

**sub**

**Add**

ALU

4

**M u x**

1

Sum

operation

MemtoReg

Instruction

Write

**Registers** Read

**Instruction memory**

register

Write data

RegWrite

data 2

**M**

**u u x x**

**M**

Zero **ALU** ALU result

|  |  |
| --- | --- |
|  | MemWrite |
| ReadAddressdata |  |
| Write **Data** |  |
| data **memory** |  |
| emRead |  |  |

offset

32 **Imm** 64

**Gen**

PCSrc

Add

4

**Shift left 1**

**Add**

**u x**

Sum

Read

**PC** address

Read register 1

Read register 2

Read data 1

ALUSrc

4 ALU operation

Zero

MemtoReg

|  |  |
| --- | --- |
|  | MemWrite |
| ReadAddressdata |  |
| Write **Data** |  |
| data **memory** |  |
|  |  |  |

Instruction

Write

**Registers** Read

**ALU**

ALU

**Instruction memory**

register

Write data

RegWrite

data 2 **M u x**

result **M**

**u**

**x**

32 **Imm** 64

**Gen**

M emRead

These three stages are fulfilled in every processor. There is a fourth stage, which is Write-Back, which is when the execution units write the result, but this is usually counted within the execution stage of the instruction cycle.

First stage of the instruction cycle: Fetch



The first stage of the instruction cycle is responsible for capturing the instructions in the [RAM](https://itigic.com/tag/ram/) memory assigned to the processor through a series of units and registers that are the following:

Program Counter or Program Counter: Which points to the next memory line where the next processor instruction is located. Its value is incremented by 1 each time a complete instruction cycle is completed or when a jump instruction changes the value of the program counter.

Memory Address Register: The MAR copies the content of the PC and sends it to the RAM through the addressing pins of the CPU, which are wired with the addressing pins of the RAM itself.

Memory Data Register or memory data register : In the event that the CPU has to perform a memory reading, what the MDR does is copy the content of that memory address to an internal register of the CPU, which is a temporary pass register before its content is copied to the Instruction Register. The MDR, unlike the MAR, is connected to the data pins of the RAM and not to the addressing pins and in the case of a write instruction the content of what you want to write in the RAM is also written in the MDR

Instruction Register: The final part of the fetch stage is the writing of the instruction in the instruction register, from which the processor control unit will copy its content for the second stage of the instruction cycle.

These 4 sub-stages occur in all processors whatever their utility, architecture and binary compatibility or what we call ISA.

Control unit



The control unit is the most complex part that exists in a processor and its tasks are as follows:

They are in charge of coordinating the movement and the order in which of the data that moves inside and outside the processor, as well as the different subunits that are in charge of it.

In general, it is considered that the units of the capture stage or Fetch are part of the hardware that we call the control unit and this hardware is also called the Front-End of a processor.

It interprets the instructions and sends them to the different execution units to which it is connected.

It is communicated to the different ALUs and execution units of the processor that act

It is responsible for capturing and decoding the instructions, but also for writing the results in the registers, caches or in the corresponding address of the RAM.



What the control unit does is decode the instructions and it does this because each instruction is actually a kind of sentence where the verb goes first and then the direct object or object on which the action is done. The subject ends up being eliminated in this internal language of computers by the fact that it is understood that it is the computer itself that executes it, so each number of bits is a sentence where the first 1 and 0 correspond to the action and the ones that come next is the data or the location of the data to be manipulated.

The second stage: Decode



There are different types of instructions and not all of them do the same, so depending on the type of instruction we need to know what execution units are going to be sent to and the most classic way of doing it is through what we call a decoder, which takes each instruction, divides it internally according to the opcode or instruction and the data or memory address where it is located.

For example in the diagram above we have the diagram of a processor with only 8 instructions, which can be encoded in only 3 bits. Each one of the instructions, once decoded, is sent to the different execution units that will resolve them.



This instruction cycle is the most complex of all and the one that defines the type of architecture. Depending on whether we have a reduced or complex set of instructions, this will affect the nature of the control unit, depending on the format of the instruction or how many are processed at the same time the decoding phase and therefore the control unit will have a different nature. other.

The easiest way to visualize what happens is to think of the instructions as trains circulating through a complex railway network and the control unit directing them to a terminal station, which is the execution unit that will be in charge of solving the instruction.

Third stage: Execute



The last stage is the execution of the instructions, in this stage the instructions are resolved, but not all types of instruction are resolved in the same way, since the way to use the hardware will depend on the function of each one of them. them, in general we have four types of instructions:

Bit movement instructions: In which the order of the bits that contain the data is manipulated.

Arithmetic instructions: Where mathematical and logical operations are carried out, these are solved in the so-called ALUs or arithmetic-logical units

Jump instructions: In which the next value of the program counter is changed, which allows the code to be used recursively.

Instructions to memory: They are with which the processor reads and writes information from the system memory.

The other point is the instruction formats, since an instruction can be applied to a data, scalar or several data at the same time, which we know as SIMD. To finish and depending on the data format, there are different types of ALUs for the execution of arithmetic instructions, for example we have integer and floating point units as differentiated units today.

Once the instruction has been completed, the result is written to a specific memory address and the next one is executed. Some instructions do not manipulate memory values but rather certain registers. Thus, the program counter register is modified by the jump instructions, if we want to read or write data then the MAR and MDR registers are manipulated.