

Is Gauss Quadrature Better than Clenshaw–Curtis?*

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Abstract. We compare the convergence behavior of Gauss quadrature with that of its younger brother, Clenshaw–Curtis. Seven-line MATLAB codes are presented that implement both methods, and experiments show that the supposed factor-of-2 advantage of Gauss quadrature is rarely realized. Theorems are given to explain this effect. First, following O’Hara and Smith in the 1960s, the phenomenon is explained as a consequence of aliasing of coefficients in Chebyshev expansions. Then another explanation is offered based on the interpretation of a quadrature formula as a rational approximation of $\log((z+1)/(z-1))$ in the complex plane. Gauss quadrature corresponds to Padé approximation at $z = \infty$. Clenshaw–Curtis quadrature corresponds to an approximation whose order of accuracy at $z = \infty$ is only half as high, but which is nevertheless equally accurate near $[-1, 1]$.

Key words. Gauss quadrature, Newton–Cotes, Clenshaw–Curtis, Chebyshev expansion, rational approximation, FFT, spectral methods

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1. Introduction. Every numerical analysis textbook has a chapter on numerical integration. Most of these present two families of $(n+1)$ -point quadrature rules: first Newton–Cotes formulas, based on equally spaced points, and then Gauss formulas, based on optimal (in a sense) points. For example, this is the pattern in the books by Henrici (1964), Isaacson and Keller (1966), Acton (1970), Dahlquist and Björck (1974), Stoer and Bulirsch (1980), Fröberg (1985), Golub and Ortega (1992), Kress (1998), and Süli and Mayers (2003) [1, 11, 22, 30, 35, 37, 39, 51, 53].

A fundamental question to ask about any family of quadrature formulas is, does it converge as $n \rightarrow \infty$, and if so, how fast?

It has been known for a long time that the Newton–Cotes formulas do not converge for a general integrand f . They converge only if f is analytic in a large region surrounding the interval of integration, and even then only in the absence of rounding errors: the problem is the $O(2^n)$ amplification associated with the Runge phenomenon [18, 46, 47]. Gauss quadrature formulas, by contrast, converge for any continuous f and have no problems with rounding errors. Moreover, they would appear to have a factor-of-2 advantage in efficiency for finite n : whereas the $(n+1)$ -point Newton–Cotes formula exactly integrates polynomials of degree n , the $(n+1)$ -point Gauss formula exactly integrates polynomials of degree $2n+1$. Gauss quadrature nodes and weights take some work to calculate, but this can be done in $O(n^2)$

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operations by solving a tridiagonal eigenvalue problem, as shown by Golub and Welsch [31].¹

However, Gauss quadrature is not the only alternative to Newton–Cotes. Here we shall compare it with *Clenshaw–Curtis quadrature*, a family of formulas based on sampling the integrand at Chebyshev points that can be implemented in $O(n \log n)$ operations using the fast Fourier transform (FFT). Clenshaw–Curtis formulas are mentioned in the numerical analysis textbooks of Johnson and Riess (1982), Ueberhuber (1997), Neumaier (2001), and Heath (2002), as well as more briefly in endnotes or exercises in Gautschi (1997) and Cheney and Kincaid (1999) [9, 27, 34, 38, 43, 57, 58]. They can also be found in monographs on numerical integration, including Davis and Rabinowitz (1967, 1975, and 1984), Brass (1977), Engels (1980), Evans (1993), Krommer and Ueberhuber (1998), and Kythe and Schäferkottter (2005) [6, 12, 18, 19, 40, 41].

Like Newton–Cotes quadrature, Clenshaw–Curtis quadrature integrates polynomials of degree n exactly. Like Gauss quadrature, it converges for all continuous f and has no problems with rounding errors. Thus on the face of it we seem to have a method that is as robust as Gauss and faster to implement for a given n , but “half as efficient.”

But here there is a surprise. In practice, Clenshaw–Curtis does not turn out to be half as efficient as Gauss. In fact, for most integrands, the two formulas are about equally accurate. This observation was made in 1968 by O’Hara and Smith [45] and reported in one or two subsequent publications, notably the books by Evans [19] and Kythe and Schäferkottter [41]. It has not become widely known, however, and has not found its way into the numerical analysis textbooks. The aim of this paper is to publicize this phenomenon and expand our understanding of it with the aid of new numerical experiments (section 3), new theorems (sections 4 and 5, especially Theorem 5.1), and a new analysis involving rational approximation in the complex plane (section 6). In a nutshell, our conclusion is that the Clenshaw–Curtis and Gauss formulas have essentially the same accuracy unless f is analytic in a sizable neighborhood of the interval of integration—in which case both methods converge so fast that the difference hardly matters. Figures 4 and 5 summarize the reason for this behavior.

Our concern throughout this article is the convergence of quadrature formulas in their “pure” form. There are many practical issues of quadrature we do not touch upon, such as indefinite integration, extrapolation, error estimation, nonuniform weight functions, stability and conditioning, compound formulas, adaptivity, endpoint and interior singularities, changes of variables, unbounded domains, and multiple dimensions. (For information about some of these matters, see, for example, [19] and [50].) We are a long way here from the design of general-purpose integration software! One reason for this emphasis is that the Gauss and Clenshaw–Curtis formulas are so basic that it is worthwhile to understand their fundamental properties in detail. Another is that even in their pure forms, quadrature formulas are used a great deal in practice as parts of other numerical computations. For example, both the Gauss and Clenshaw–Curtis formulas are employed in the numerical solution of ODEs and PDEs by spectral methods [7, 55]. Another example, the one that got me involved in this topic, is the chebfun system in object-oriented MATLAB, which takes advantage of the power and speed of Clenshaw–Curtis quadrature to integrate functions on the fly that may be defined by as many as millions of data points [4].

All in all, it would seem that the Gauss and Clenshaw–Curtis formulas should perhaps be regarded as equally valuable and fundamental, with the former having an edge in elegance and the latter in simplicity.

¹An indispensable source of information on Gauss quadrature is the masterful 1981 survey by Gautschi [26].

2. Gauss and Clenshaw–Curtis Formulas. Throughout this article we are concerned with what are known as *interpolatory quadrature formulas*. A continuous function f on $[-1, 1]$ is given (transplantation to a general interval $[a, b]$ is trivial), and we seek to approximate the integral

$$(2.1) \quad I = I(f) = \int_{-1}^1 f(x) dx$$

by sums

$$(2.2) \quad I_n = I_n(f) = \sum_{k=0}^n w_k f(x_k)$$

for various n , where the *nodes* $\{x_k\}$ depend on n but not f . The *weights* $\{w_k\}$ are defined uniquely by the property that I_n is equal to the integral of the degree $\leq n$ polynomial interpolant through the data points. Thus, by construction, (2.2) gets the exactly correct answer if f is a polynomial of degree $\leq n$.

Newton–Cotes formulas are defined by taking the nodes to be equally spaced from -1 to 1 . The Newton–Cotes formulas of low order are important as building blocks for discretizations of differential equations and other numerical methods. Their properties as $n \rightarrow \infty$ are very bad, however, and we shall not discuss them further except for brief reappearances in Figures 1 and 4.

Gauss formulas are defined by choosing the nodes optimally in the sense of maximizing the degree of polynomials that (2.2) integrates exactly. Since there are $n + 1$ nodes, it is not surprising that the attainable degree is $n + 1$ orders higher than in the generic case, i.e., degree $2n + 1$. Gauss quadrature formulas were discovered by Gauss in his mid-thirties [25], and it is a measure of the extraordinary productivity of Gauss’s career that this great discovery merits hardly a mention in many accounts of his life. For these formulas are truly powerful. Their impact was mainly theoretical before the advent of computers, because of the difficulty of determining the nodes and weights and also of evaluating integrands at irrational arguments, but in the computer era, especially since the appearance of a paper by Golub and Welsch in 1969 (following earlier work of Goertzel in 1954, Wilf in 1962, and Gordon in 1968), it has been practical as well [31]. We shall not give details but just offer the following MATLAB function `gauss.m`, which is almost the same as the code by the same name in [55]:

```
function I = gauss(f,n)           % (n+1)-pt Gauss quadrature of f
beta = .5./sqrt(1-(2*(1:n)).^(-2)); % 3-term recurrence coeffs
T = diag(beta,1) + diag(beta,-1); % Jacobi matrix
[V,D] = eig(T);                  % eigenvalue decomposition
x = diag(D); [x,i] = sort(x);    % nodes (= Legendre points)
w = 2*V(1,i).^2;                 % weights
I = w*feval(f,x);                % the integral
```

For example, the command `gauss(@cos,6)` yields 1.68294196961579, which is correct to full precision. The command `gauss(@cos,1000)` yields the same result, though it takes 17 seconds on my laptop.²

²Here are Stroud and Secrest in 1966, working on a Control Data 1604: “The formula for $n = 512$ took about 4.5 hours to calculate and 0.5 hour to check” [52]. Thus in forty years we have a speedup from hours on a mainframe to seconds on a laptop. This comparison is complicated, however, by the fact that the two calculations differed in many ways. For example, Stroud and Secrest worked in multiple precision and before the days of Golub and Welsch [31]; meanwhile `gauss.m` uses the Golub–Welsch algorithm in standard precision but in a form requiring $O(n^3)$ rather than $O(n^2)$ operations, since the `eig` command does not take advantage of the fact that the matrix is tridiagonal.

The idea of Clenshaw–Curtis quadrature is to use Chebyshev points instead of optimal nodes. There are three main variations on this theme (see [44] for others):

- Chebyshev roots in $(-1, 1)$: Fejér’s “first rule” (1933) [20];
- Chebyshev extrema in $(-1, 1)$: Fejér’s “second rule” (1933) [20, 21];
- Chebyshev extrema in $[-1, 1]$: Clenshaw and Curtis (1960) [10].

The first and the third variants are also called “classical” and “practical” Clenshaw–Curtis formulas, respectively. (A number of the relevant papers published, especially in the 1960s, make no reference to Fejér.) It seems that the last choice may be the most accurate [16] as well as having the easiest connection with the FFT,³ and also the potential advantage of convenient reuse of nodes when n is doubled; it is certainly the one used mainly in practice. The nodes in question, which we shall simply call *Chebyshev points*, are defined by

$$(2.3) \quad x_j = \cos \frac{j\pi}{n}, \quad 0 \leq j \leq n.$$

The *Clenshaw–Curtis quadrature formula* is the formula (2.2) based on these nodes. A better name might have been “Chebyshev” or “Fejér”—indeed, Clenshaw and Curtis call it “the Chebyshev formula”—but the term “Clenshaw–Curtis” is standard.

Clenshaw and Curtis published their paper in 1960, before the introduction of the FFT in 1965. Soon afterwards, the connection with the FFT was pointed out by Gentleman [28, 29]. Again we shall not give details but offer a MATLAB code with the same functionality as `gauss.m`:

```
function I = clenshaw_curtis(f,n)           % (n+1)-pt C-C quadrature of f
x = cos(pi*(0:n)'/n);                     % Chebyshev points
fx = feval(f,x)/(2*n);                     % f evaluated at these points
g = real(fft(fx([1:n+1 n:-1:2])));         % Fast Fourier Transform
a = [g(1); g(2:n)+g(2*n:-1:n+2); g(n+1)]; % Chebyshev coefficients
w = 0*a'; w(1:2:end) = 2./(1-(0:2:n).^2); % weight vector
I = w*a;                                   % the integral
```

To get full accuracy for the cosine example, we now need `clenshaw_curtis(@cos,11)`. Increasing n from 11 to 10^3 or 10^6 gives the same result. The former runs in less than a millisecond on my laptop, and the latter takes 1.6 seconds.

From the point of view of integrating polynomials exactly, Clenshaw–Curtis nodes appear like Newton–Cotes nodes: both are exact merely for degree n . Yet it is clear from Figure 1 that in a literal sense, at least, they are closer to Gauss nodes. Certainly they have the same asymptotic distribution as $n \rightarrow \infty$, the density $n\pi^{-1}(1-x^2)^{-1/2}$ that is well known to be optimal in various senses for polynomial approximation on $[-1, 1]$ [55, Chap. 5].⁴

3. Experimental Comparison. Figure 2 shows the convergence as $n \rightarrow \infty$ of the Gauss and Clenshaw–Curtis formulas for six functions $f(x)$. (This experiment is adapted from Outputs 30b and 30c of [55].) We begin with the monomial $f(x) = x^{20}$, the algebraic analogue of what in the periodic context would be called a band-limited

³For FFT evaluations of the other two rules, see [59].

⁴The clustering of optimal grids for polynomial interpolation implies, however, that in certain cases alternative nonpolynomial methods can be more efficient by a factor as high as $\pi/2$. One method for recovering this factor is described in [2]; we propose another in [33]. For some of the further implications of this “excessive” clustering of nodes near endpoints, see [14].

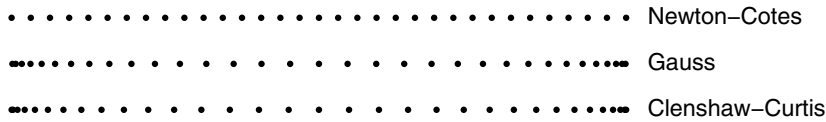


Fig. 1 Newton–Cotes, Gauss, and Clenshaw–Curtis quadrature points in $[-1, 1]$ for $n = 32$. The latter two sets have asymptotically the same clustering near ± 1 as $n \rightarrow \infty$.

function. Here the factor of 2 is in one sense clearly visible: the Gauss formula is exact for $n \geq 10$, the Clenshaw–Curtis formula for $n \geq 20$. Even so, it is apparent that for smaller values of n , the Gauss advantage in efficiency (i.e., size of n needed to achieve a certain accuracy) is less than a factor of 2. The second function is e^x , which is entire, i.e., analytic throughout the complex plane. Here again Gauss appears more efficient than Clenshaw–Curtis, but by less than a factor of 2.

Given $f \in C[-1, 1]$ and $n \geq 0$, let p_n^* be the unique best approximation to f on $[-1, 1]$ of degree $\leq n$ with respect to the supremum norm $\|\cdot\| = \|\cdot\|_\infty$, and define $E_n^* = \|f - p_n^*\|$. The solid curves in each panel of Figure 2 show the errors E_n^* and E_{2n+1}^* . (We computed these numbers via Carathéodory–Fejér approximation [32] using the chebfun system [4], a convenient environment for all kinds of explorations related to this paper.) We shall analyze these curves in the next section, but, for now, we mention just the essential point: Clenshaw–Curtis quadrature converges faster than E_n^* .

The next four functions are successively less smooth. First we have e^{-x^2} , which is again entire, but now with stronger growth $O(e^{|x|^2})$ as $x \rightarrow \infty$ in the complex plane. Here again we see a difference between Gauss and Clenshaw–Curtis quadrature, but with an efficiency ratio less than 2. Next comes $f(x) = (1 + 16x^2)^{-1}$, which is analytic in a neighborhood of $[-1, 1]$ but not throughout the complex plane, and now we see that the two quadrature formulas are quite close together. (This example is discussed further in section 6.) For the last two functions, there is again little difference between Gauss and Clenshaw–Curtis quadrature. We see this for both the C^∞ function e^{-1/x^2} and the nonsmooth function $|x|^3$.

These and similar experiments suggest the following conclusions:

- As expected, Gauss quadrature matches best degree $2n + 1$ polynomial approximation: $|I - I_n| \approx E_{2n+1}^*$.
- However, Clenshaw–Curtis quadrature does much better than $|I - I_n| \approx E_n^*$. In fact, for functions that are not analytic in a sizable neighborhood of $[-1, 1]$, it too comes close to $|I - I_n| \approx E_{2n+1}^*$.
- Thus Gauss quadrature significantly outperforms Clenshaw–Curtis quadrature only for functions analytic in a sizable neighborhood of $[-1, 1]$.
- For such functions, the convergence of both methods is very fast. Thus Clenshaw–Curtis essentially never requires many more function evaluations than Gauss to converge to a prescribed accuracy.

We are not the first to observe these surprises. Clenshaw and Curtis themselves recorded the same effect [10]:

It may be of interest to note that a program based on the proposed method has been prepared for the DEUCE at the National Physical Laboratory. The upper limit M of the number of terms is 64, and the program evaluated the integral $\int_{-1}^{+1} |x + \frac{1}{2}|^{\frac{1}{2}} dx$ with an error of 0.00078. This may be compared

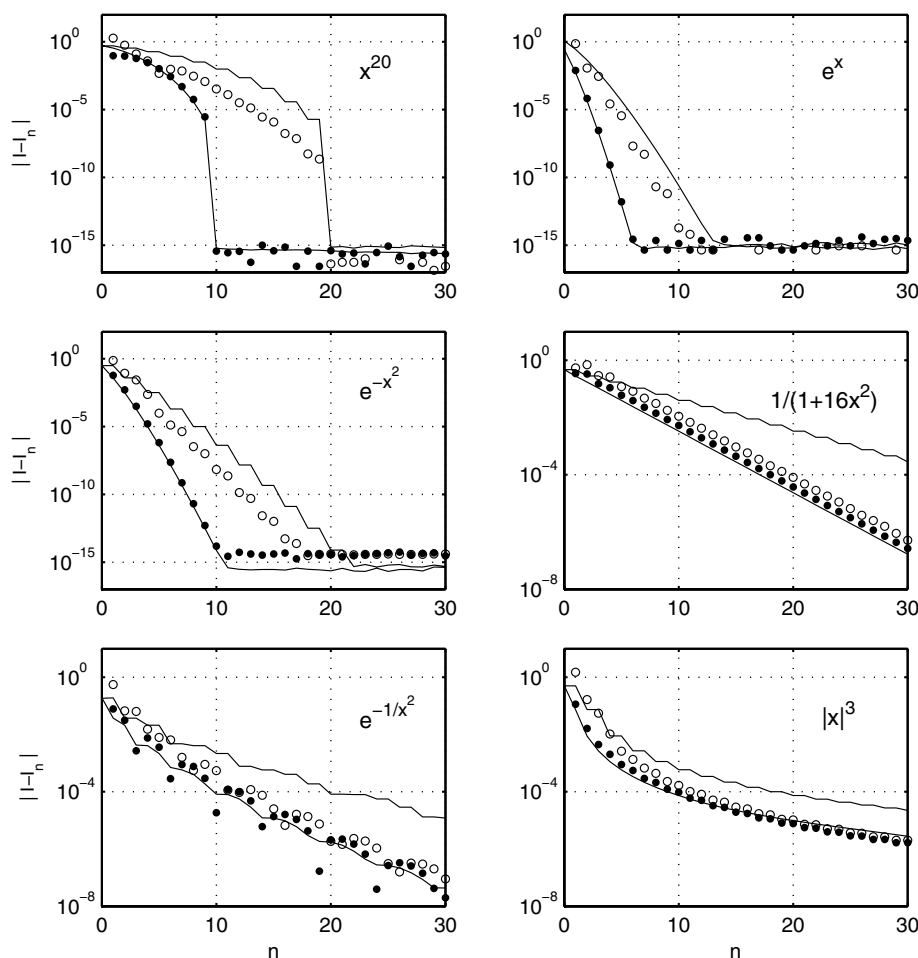


Fig. 2 Convergence in floating-point arithmetic of Gauss (dots) and Clenshaw–Curtis (circles) quadrature for six functions f on $[-1, 1]$ for n ranging from 1 to 30. The first function is polynomial, the second and third entire, the fourth analytic, the fifth C^∞ , and the sixth C^2 . The solid curves show the minimax approximation errors E_n^* (upper) and E_{2n+1}^* (lower).

with the error of 0.00317 committed by the 32-point Gauss formula, and of 0.00036 by the 64-point Gauss formula. We see that the Chebyshev formula, which is much more convenient than the Gauss, may sometimes nevertheless be of comparable accuracy.

Figure 3 examines this function considered by Clenshaw and Curtis for a collection of values of n . There is no advantage whatever for Gauss quadrature. (The reason the plot does not match the numbers they reported is that it shows data for $n = 64$, whereas their numbers correspond to $n = 63$.)

A more comprehensive claim of the practical equivalence of Clenshaw–Curtis and Gauss quadrature appeared in a 1968 paper of O’Hara and Smith [45], building on related work by Elliott [16]. They wrote

... the Clenshaw–Curtis method gives results nearly as accurate as Gaussian quadratures for the same number of abscissae ...

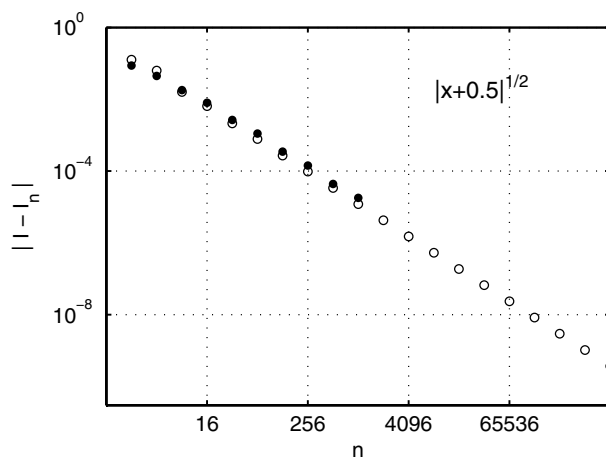


Fig. 3 Convergence of Gauss (dots) and Clenshaw–Curtis (circles) quadrature for the example $f(x) = |x + \frac{1}{2}|^{1/2}$ of Clenshaw and Curtis, shown on log-log axes. The Gauss formula is applied only up to $n = 2^{10}$, since larger values take too much time and memory. Of course, the point of this experiment is to elucidate fundamental convergence properties; there are many ways in which a function this simple could be integrated numerically with far fewer sample points.

and gave a justification of this statement based on magnitudes of Chebyshev coefficients. Their conclusions were highlighted subsequently in the book of Evans [19]. We turn to such matters in the next two sections, and for the O’Hara and Smith argument in particular, see (5.5) and the discussion just afterwards. For a variety of convergence estimates for Gauss, Clenshaw–Curtis, and related quadrature formulas, a particularly valuable reference is the book by Brass (in German) [6].

4. Quadrature, Polynomial Approximation, and Chebyshev Expansions. How can we explain the surprising accuracy of Clenshaw–Curtis quadrature? To prepare the way, we begin with a discussion of some connections between quadrature formulas, polynomial approximations, and Chebyshev expansions. Though the flavor of these developments is standard, some of the estimates we give are sharper than those one normally finds in the literature; this applies to (4.6), (4.10), (4.12), and (4.13), as well as the Jackson theorem of footnote 6. (This does not mean they are as sharp as possible! In particular, G. Mastroianni has pointed out to me that stronger estimates can be obtained based on results in [13] and [14].) We consider Chebyshev expansions for two reasons: first, because they provide near-best polynomial approximations and thus near-sharp bounds on Gauss quadrature errors; second, because they have a special connection with Clenshaw–Curtis quadrature.

Given a continuous function f on $[-1, 1]$, let the degree $\leq n$ polynomial best approximation p_n^* and error E_n^* be defined as in the last section. The rationale for including the solid curves in Figure 2 is the following result due to Bernstein [5].

THEOREM 4.1. *If the weights w_k in (2.2) are nonnegative, as is true for both Gauss and Clenshaw–Curtis quadrature, then, for any $f \in C[-1, 1]$ and $n \geq 0$,*

$$(4.1) \quad |I - I_n| \leq 4E_n^*,$$

and $I_n \rightarrow I$ as $n \rightarrow \infty$. In the special case of Gauss quadrature, (4.1) can be improved to

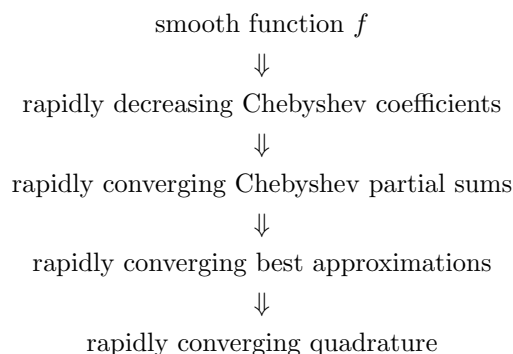
$$(4.2) \quad |I - I_n| \leq 4E_{2n+1}^*.$$

Proof. We take advantage of the fact that since (2.2) integrates polynomials of degree $\leq n$ exactly, $(I - I_n)(f) = (I - I_n)(f - p_n^*)$. From (2.2), using the fact that the interval $[-1, 1]$ has length 2, we thus have

$$(4.3) \quad \begin{aligned} |I(f) - I_n(f)| &= |I(f - p_n^*) - I_n(f - p_n^*)| \\ &\leq |I(f - p_n^*)| + |I_n(f - p_n^*)| \leq 2E_n^* + \sum_{k=0}^n |w_k| E_n^*. \end{aligned}$$

Since the quadrature formula is interpolatory, $\sum w_k = 2$, and if the weights are nonnegative, this implies $\sum |w_k| = 2$, from which (4.1) follows, and similarly (4.2) in the case of Gauss quadrature since in that case p_n^* can be replaced by p_{2n+1}^* . The positivity of the weights was proved for Gauss quadrature by Stieltjes in 1884 and for Clenshaw–Curtis quadrature by Imhof in 1963 [36]. The convergence of I_n to I is a consequence of (4.1) together with the Weierstrass approximation theorem. $\square$

Theorem 4.1 tells us that if the best approximants to f converge rapidly as $n \rightarrow \infty$, then I_n will converge rapidly to I . The next step is to combine this theorem with results of approximation theory to the effect that if f is smooth, its best approximants converge rapidly. We shall consider some results of this kind derived from the *Chebyshev series* for a function $f \in C[-1, 1]$. The logic of the development can be summarized like this:



However, we shall see in the next section that sharper results can sometimes be obtained by bypassing steps 3 and 4 and going directly from the Chebyshev coefficients to the quadrature formulas.

Let T_j denote the Chebyshev polynomial of degree j , $T_j(x) = \cos(j \cos^{-1} x)$, which equioscillates between $j + 1$ extrema ± 1 on $[-1, 1]$. The Chebyshev series for $f \in C[-1, 1]$ is defined by [49, p. 165]

$$(4.4) \quad f(x) = \sum_{j=0}^{\infty}{}' a_j T_j(x), \quad a_j = \frac{2}{\pi} \int_{-1}^1 \frac{f(x) T_j(x)}{\sqrt{1-x^2}} dx,$$

where the prime indicates that the term with $j = 0$ is multiplied by $1/2$. (These formulas are nothing more than the transplantation to $x = \cos \theta$ of the Fourier series

for the 2π -periodic even function $f(\cos \theta)$.) The equals sign in the first formula is justified under the mild condition that f is Dini-continuous, in which case the series converges uniformly to f ; this is true in particular if f is Hölder continuous for any Hölder exponent > 0 [49, p. 168].

Our first step is to show that if f is smooth, its Chebyshev coefficients decrease rapidly. We consider two smoothness conditions: a k th derivative satisfying a condition related to bounded variation, or analyticity in a neighborhood of $[-1, 1]$. Regarding the former, let $\|\cdot\|_T$ be the Chebyshev-weighted 1-norm defined by

$$(4.5) \quad \|u\|_T = \left\| \frac{u'(x)}{\sqrt{1-x^2}} \right\|_1.$$

This norm is defined via a Stieltjes integral for any u of bounded variation (see, e.g., [61, p. 220]), though $\|u\|_T$ may be infinite depending on the behavior near $x = \pm 1$. The condition of interest for our theorems is that $\|f^{(k)}\|_T$ should be finite, where $f^{(k)}$ is the k th derivative of f .⁵

THEOREM 4.2. *If $f, f', \dots, f^{(k-1)}$ are absolutely continuous on $[-1, 1]$ and if $\|f^{(k)}\|_T = V < \infty$ for some $k \geq 0$, then for each $n \geq k + 1$,*

$$(4.6) \quad |a_n| \leq \frac{2V}{\pi n(n-1) \cdots (n-k)}.$$

If f is analytic with $|f(z)| \leq M$ in the region bounded by the ellipse with foci ± 1 and major and minor semiaxis lengths summing to $\rho > 1$, then for each $n \geq 0$,

$$(4.7) \quad |a_n| \leq \frac{2M}{\rho^n}.$$

Proof. For (4.7), see [49, p. 175]. As for (4.6), I do not know if this result can be found in print; I am grateful to Endre Süli for suggesting the following argument. We proceed by transplanting to the $\theta = \cos^{-1} x$ variable and integrating by parts $k + 1$ times; the assumptions ensure that the first k derivatives exist in the ordinary sense and that $f^{(k)}$ can in turn be written as a Stieltjes integral. Since $dx = -\sin \theta d\theta$ and $\sqrt{1-x^2} = \sin \theta$, the first integration by parts can be written as

$$\begin{aligned} a_n &= \frac{2}{\pi} \int_{-1}^1 \frac{f(x) T_n(x)}{\sqrt{1-x^2}} dx \\ &= \frac{2}{\pi} \int_0^\pi f(\cos \theta) \cos(n\theta) d\theta \\ &= \frac{2}{\pi n} \int_0^\pi f'(\cos \theta) \sin(n\theta) \sin \theta d\theta; \end{aligned}$$

the boundary terms vanish because of the sine. The trigonometric identity $\sin \theta_1 \sin \theta_2 = \frac{1}{2} \cos(\theta_1 - \theta_2) - \frac{1}{2} \cos(\theta_1 + \theta_2)$ converts this to

$$a_n = \frac{2}{\pi n} \int_0^\pi f'(\cos \theta) \left[\frac{\cos((n-1)\theta)}{2} - \frac{\cos((n+1)\theta)}{2} \right] d\theta,$$

⁵Our attention to this condition derives from a moral principle: the exponent should come out right for the function $f(x) = |x|$. For this function we can take $k = 1$ in Theorem 4.2, giving $a_n = O(n^{-2})$.

which implies

$$\begin{aligned} |a_n| &\leq \frac{2}{\pi n} \|f'(\cos \theta)\|_1 \left\| \frac{\cos((n-1)\theta)}{2} - \frac{\cos((n+1)\theta)}{2} \right\|_\infty \\ &= \frac{2}{\pi n} \|f(x)\|_T \left\| \frac{\cos((n-1)\theta)}{2} - \frac{\cos((n+1)\theta)}{2} \right\|_\infty \end{aligned}$$

since $d\theta = dx/\sqrt{1-x^2}$. The L^∞ norm is bounded by 1, and thus we have established $|a_n| \leq 2V^{(0)}/\pi n$ with $V^{(0)} = \|f\|_T$. Further integrations by parts bring in higher derivatives of f and corresponding higher variations up to $V^{(k)} = V = \|f^{(k)}\|_T$. More and more cosine terms appear, but the coefficients are such that their sum always has L^∞ norm at most 1. Just as the first integration by parts introduced a factor n in the denominator, the second leads to factors $n-1$ and $n+1$, the third to factors $n-2$, n , and $n+2$, and so on. To keep the formulas simple we do not keep track of all these different denominators but weaken the inequality slightly by replacing them all with $n-1$ at the second differentiation, $n-2$ at the third, and so on up to $n-k$ at the $(k+1)$ st differentiation. The result is (4.6). $\square$

Our next step is to consider polynomial approximations obtained as partial sums of the Chebyshev series. For any $n \geq 0$, define

$$(4.8) \quad f_n(x) = \sum_{j=0}^n a_j T_j(x)$$

and

$$(4.9) \quad E_n^T = \|f - f_n\| \leq \sum_{j=n+1}^{\infty} |a_j|.$$

(The polynomial f_n can be interpreted as the least-squares projection of f onto the set of polynomials of degree $\leq n$ with respect to the weight $(1-x^2)^{-1/2}$. It can also be obtained from f by convolution in the θ variable with a certain n -dependent kernel [49, p. 166].) The following theorem is a corollary of Theorem 4.2 and (4.9).

THEOREM 4.3. *If $f, f', \dots, f^{(k-1)}$ are absolutely continuous on $[-1, 1]$ and if $\|f^{(k)}\|_T = V < \infty$ for some $k \geq 1$, then for each $n \geq k+1$,*

$$(4.10) \quad E_n^T \leq \frac{2V}{\pi k(n-k)^k}.$$

If f is analytic with $|f(z)| \leq M$ in the region bounded by the ellipse with foci ± 1 and major and minor semiaxis lengths summing to $\rho > 1$, then for each $n \geq 0$,

$$(4.11) \quad E_n^T \leq \frac{2M}{(\rho-1)\rho^n}.$$

Proof. Equation (4.11) follows readily from (4.7) and (4.9). (This bound is erroneously reported without the factor of 2 on p. 133 of [42].) Similarly, (4.10) follows from (4.6) and (4.9) if we use the inequality

$$\sum_{j=n+1}^{\infty} \frac{1}{j(j-1)\cdots(j-k)} \leq \sum_{j=n+1}^{\infty} \frac{1}{(j-k)^{k+1}} \leq \int_n^{\infty} \frac{dx}{(x-k)^{k+1}} = \frac{1}{k(n-k)^k}. \quad \square$$

Now how does E_n^T , the error of the truncated Chebyshev series, compare with the best approximation error E_n^* ? In one direction the answer is obvious, since the optimality of E_n^* implies $E_n^* \leq E_n^T$.⁶ In the other direction it is well known that the gap is never bigger than $O(\log n)$. The following theorem states this result together with another lower bound on E_n^* .

THEOREM 4.4. *For any $f \in C[-1, 1]$ and $n \geq 0$,*

$$(4.12) \quad \frac{\pi}{4}|a_{n+1}| \leq E_n^* \leq E_n^T \leq \left(2 + \frac{4}{\pi^2} \log(2n+1)\right) E_n^*.$$

In particular, $E_n^T < 4.2E_n^$ for $n \leq 100$ and $E_n^T < 7.9E_n^*$ for $n \leq 10^6$.*

Proof. The first inequality is due to Bernstein; see [8, p. 131] or [49, p. 171]. The last inequality follows from the fact that the Lebesgue constant for Chebyshev projection, which is the same as the Lebesgue constant for Fourier projection defined originally by Lebesgue (1909), is bounded by $1 + (4/\pi^2) \log(2n+1)$ [23]. The $\sim (4/\pi^2) \log n$ dependence was known already to Fejér (1910) and Bernstein (1912). $\square$

By combining Theorems 4.1 and 4.3, we get bounds on $I - I_n$, which Theorem 4.4 indicates are likely to be reasonably sharp. Here are bounds for Gauss quadrature. The bound (4.14) is due to Rabinowitz, appearing as equation (18) of [48] and also as Theorem 90 of [6]. Bounds based on the assumption of analyticity in an ellipse were first derived by Bernstein [5]; for a discussion, see [26, p. 114].

THEOREM 4.5. *Let Gauss quadrature be applied to a function $f \in C[-1, 1]$. If $f, f', \dots, f^{(k-1)}$ are absolutely continuous on $[-1, 1]$ and $\|f^{(k)}\|_T = V < \infty$ for some $k \geq 1$, then for each $n \geq k/2$,*

$$(4.13) \quad |I - I_n| \leq \frac{32V}{15\pi k(2n+1-k)^k}.$$

If f is analytic with $|f(z)| \leq M$ in the region bounded by the ellipse with foci ± 1 and major and minor semiaxis lengths summing to $\rho > 1$, then for each $n \geq 0$,

$$(4.14) \quad |I - I_n| \leq \frac{64M}{15(1-\rho^{-2})\rho^{2n+2}}.$$

Partial proof. Suppose the factors $32/15$ and $64/15$ in these statements are increased to 8 and the factor $1 - \rho^{-2}$ in the denominator of (4.14) is decreased to $1 - \rho$. Then (4.13) follows from (4.2) and (4.10), and (4.14) from (4.2) and (4.11). The sharpening of these factors to complete the proof requires consideration of individual Chebyshev coefficients and is discussed in the next section. $\square$

The bounds of Theorem 4.5 are valid for Clenshaw–Curtis quadrature, except with $2n+1$ replaced by n . However, we do not write these inequalities down, as we have seen that they are highly pessimistic and our aim is to explain that effect. Ideally, we would like to establish bounds for Clenshaw–Curtis quadrature that are not much different from those of Theorem 4.5—i.e., with “ $2n$ ”, not “ n ”. The next section will present such an extension to Clenshaw–Curtis quadrature of (4.13), and section 6 will point the way toward a possible extension of (4.14).

⁶Thus (4.10) provides a proof of a variant of one of the *Jackson theorems* of approximation theory [8, p. 147]: $E_n^* \leq 2V/\pi k(n-k)^k$.

5. Explanation I: Chebyshev Expansions and Aliasing. Let us immediately state our conclusion: the first half of Theorem 4.5 holds for Clenshaw–Curtis as well as Gauss quadrature. The sole change needed is to replace the condition $n \geq k/2$ by the condition that n is sufficiently large. Thus Clenshaw–Curtis has essentially the same algebraic rate of convergence as Gauss quadrature for several-times-differentiable functions.

THEOREM 5.1. *Let Clenshaw–Curtis quadrature be applied to a function $f \in C[-1, 1]$. If $f, f', \dots, f^{(k-1)}$ are absolutely continuous on $[-1, 1]$ and $\|f^{(k)}\|_T = V < \infty$ for some $k \geq 1$, then for all sufficiently large n ,*

$$(5.1) \quad |I - I_n| \leq \frac{32V}{15\pi k(2n+1-k)^k}.$$

“Sufficiently large n ” means $n > n_k$ for some n_k that depends on k but not f or V .

Notice the key implication: the factor 2^{-k} in the error bound (4.13) for Gauss quadrature, which one thinks of as coming from its doubled order of accuracy, applies to Clenshaw–Curtis quadrature too. In this section we shall derive this bound. The theorem is new, but it is built on the ideas of Elliott [16] and O’Hara and Smith [45] from forty years ago, combined with the “rough and ready” estimates of Rabinowitz [48].

The crucial fact is that of aliasing. Elementary algebra shows that on the grid in $[0, 2\pi]$ of $2n$ equally spaced points $\theta_j = \pi j/n$, $0 \leq j \leq 2n-1$, the functions $\cos((n+p)\pi\theta_j)$ and $\cos((n-p)\pi\theta_j)$ are indistinguishable for any integer p . We say that the wave numbers $(n+p)\pi$ and $(n-p)\pi$ are *aliases* of one another on this grid. Transplanting to the variable $x = \cos \theta$ gives the following result.

THEOREM 5.2. *For any integer p with $0 \leq p \leq n$,*

$$(5.2) \quad T_{n+p}(x_j) = T_{n-p}(x_j) \quad (0 \leq j \leq n)$$

on the Chebyshev grid (2.3). Consequently Clenshaw–Curtis quadrature gives the same result for both functions,

$$(5.3) \quad I_n(T_{n+p}) = I_n(T_{n-p}) = I(T_{n-p}) = \begin{cases} 0 & \text{if } n \pm p \text{ is odd,} \\ \frac{2}{1 - (n-p)^2} & \text{if } n \pm p \text{ is even,} \end{cases}$$

and the error in integrating T_{n+p} is

$$(5.4) \quad I(T_{n+p}) - I_n(T_{n+p}) = \begin{cases} 0 & \text{if } n \pm p \text{ is odd,} \\ \frac{8pn}{n^4 - 2(p^2 + 1)n^2 + (p^2 - 1)^2} & \text{if } n \pm p \text{ is even.} \end{cases}$$

Proof. The equality (5.2) was established in the discussion above it. Since Clenshaw–Curtis quadrature is exact for polynomials of degree $\leq n$, the identity (5.3) follows from the standard formula $\int_{-1}^1 T_j(x) dx = 2/(1-j^2)$ for the integral of a Chebyshev polynomial of even order (which appeared in the penultimate line of `clenshaw_curtis` in section 2). Finally, taking the difference of the integrals with $j = n+p$ and $j = n-p$ gives (5.4). $\square$

From (5.4) we can see why Clenshaw–Curtis quadrature is unexpectedly accurate. Suppose n is even for simplicity. Then the first few terms in the Chebyshev expansion of f that contribute to the error $I(f) - I_n(f)$ are

$$a_{n+2}T_{n+2}, a_{n+4}T_{n+4}, \dots$$

However, the exact integrals of these Chebyshev polynomials are small, of order $O(n^{-2})$, and, thanks to aliasing, their values as computed by Clenshaw–Curtis quadrature will also be small. From (5.4), we see that if n is large, the errors associated with the terms above are approximately

$$(5.5) \quad \frac{16}{n^3}|a_{n+2}|, \quad \frac{32}{n^3}|a_{n+4}|, \quad \dots$$

The numbers grow as the subscript increases from n towards $2n$, but even for a “typical” term a_{n+p} , the error will be of order n^{-2} . For a numerical illustration, suppose $n = 50$. Then the $(n+1)$ -point Clenshaw–Curtis formula integrates T_{52} inexactly, but the error is only about 0.0001. Similarly its errors for T_{60} , T_{70} , T_{80} , and T_{90} are about 0.0006, 0.002, 0.006, and 0.02, respectively. By contrast the Gauss formula with $n = 50$ integrates all the polynomials up to T_{101} exactly, but after that its error jumps immediately to order 1: for T_{102} , about -1.6 .

We can now derive (5.1) and, in the process, explain how the constants 8 and $1 - \rho$ in our partial proof of Theorem 4.5 can be improved to $32/15$ and $64/15$ and $1 - \rho^{-2}$.

Proof of Theorem 5.1 and completion of proof of Theorem 4.5. From the Chebyshev expansion of f we know that the quadrature error can be written as

$$I(f) - I_n(f) = \sum_{j=0}^{\infty} a_j (I(T_j) - I_n(T_j)).$$

This implies

$$|I(f) - I_n(f)| \leq S_1 + S_2 + S_3 + S_4,$$

where

$$\begin{aligned} S_1 &= \sum_{j=0}^n |a_j| |I(T_j) - I_n(T_j)|, \\ S_2 &= \sum_{j=n+1}^{2n-\lfloor n^{1/3} \rfloor} |a_j| |I(T_j) - I_n(T_j)|, \\ S_3 &= \sum_{j=2n+1-\lfloor n^{1/3} \rfloor}^{2n+1} |a_j| |I(T_j) - I_n(T_j)|, \\ S_4 &= \sum_{j=2n+2}^{\infty} |a_j| |I(T_j) - I_n(T_j)|. \end{aligned}$$

The term S_1 is zero since the quadrature formula is interpolatory. First let us consider Gauss quadrature. In this case S_2 and S_3 are zero too, so our only task is to estimate S_4 . In the last section, we implicitly used the inequality $|I(T_j) - I_n(T_j)| \leq 4$ for each j . However, since $j \geq 4$ and T_j is odd when j is odd, from (5.3) it can be seen that we actually have

$$(5.6) \quad |I(T_j) - I_n(T_j)| \leq \begin{cases} 32/15 & \text{if } j \text{ is even,} \\ 0 & \text{if } j \text{ is odd,} \end{cases}$$

since $2 + 2/(4^2 - 1) = 32/15$. Equations (4.13) and (4.14) follow by combining this inequality with (4.6) and (4.7), respectively, and, in the case of (4.13), using an estimate of a sum analogous to that in the proof of Theorem 4.3.

Next let us consider Clenshaw–Curtis quadrature. From (5.4) it can be seen that the factors $|I(T_j) - I_n(T_j)|$ that appear in S_2 are all of order at worst $n^{-2/3}$ and, by Theorem 4.2, the coefficients a_j are of order Vn^{-k-1} . Thus S_2 consists of $O(n)$ terms of size $O(Vn^{-k-5/3})$, for a total magnitude $O(Vn^{-k-2/3})$. The number S_3 consists of $O(n^{1/3})$ terms of size $O(Vn^{-k-1})$, for a total magnitude again $O(Vn^{-k-2/3})$. Finally, consider S_4 once more. The statement of Theorem 5.1 allows us the freedom to increase n a little further. In particular, taking $n \geq 2$ is enough to ensure $j \geq 6$, in which case (5.6) holds with the constant $32/15$ improved to $2 + 2/(6^2 - 1) = 72/35$. Combining this with (4.6) and putting the pieces together, we find

$$S_1 + S_2 + S_3 + S_4 \leq O(Vn^{-k-2/3}) + \frac{72V/35}{\pi k(2n+1-k)^k} < \frac{32V/15}{\pi k(2n+1-k)^k}$$

for large enough n . $\square$

6. Explanation 2: Rational Approximation of $\log((z+1)/(z-1))$. The development so far has not followed the sequence that led to this article being written. Some time ago I undertook to investigate analogues for $[-1, 1]$ of observations put forward in [56] about quadrature and rational approximation on $(-\infty, 0]$. As soon as images like Figures 4 and 5 appeared, I knew there was a story to be told about Clenshaw–Curtis quadrature.

To set the stage, we turn to an old idea for estimating the accuracy of quadrature formulas applied to analytic functions: Cauchy integrals in the complex plane. As always, we are interested in the relationship of the integral I of (2.1) to its discrete approximation I_n of (2.2) for a function $f \in C[-1, 1]$. Suppose that f is analytic in a neighborhood of $[-1, 1]$, and let Γ be a contour in the region of analyticity that winds once around $[-1, 1]$ in the counterclockwise direction. Then using Cauchy's integral formula we can write

$$(6.1) \quad I = \int_{-1}^1 f(x) dx = \int_{-1}^1 \frac{1}{2\pi i} \int_{\Gamma} \frac{f(z)}{z-x} dz dx = \frac{1}{2\pi i} \int_{\Gamma} f(z) \phi(z) dz,$$

where $\phi(z)$ is given by

$$(6.2) \quad \phi(z) = \int_{-1}^1 \frac{dx}{z-x} = \log \frac{z+1}{z-1}.$$

To be precise, ϕ is defined as the analytic function in $\mathbb{C} \cup \{\infty\} \setminus [-1, 1]$ given by the branch of $\log((z+1)/(z-1))$ that is positive for $z \in (1, \infty)$. At the same time, using complex residue calculus, we can also express I_n as a contour integral over Γ :

$$(6.3) \quad I_n = \sum_{k=0}^n w_k f(x_k) = \frac{1}{2\pi i} \int_{\Gamma} f(z) r_n(z) dz,$$

where r_n is defined by

$$(6.4) \quad r_n(z) = \sum_{k=0}^n \frac{w_k}{z-x_k}.$$

Combining (6.1) and (6.3) now gives a beautiful result, which Takahasi and Mori stated in [54] though they were possibly not the first to do so. For more on quadrature error estimates from contour integrals, see [15].

THEOREM 6.1. *Let f be analytic in a neighborhood of $[-1, 1]$ and let Γ be a contour contained in the region of analyticity of f that encloses $[-1, 1]$ once in the counterclockwise direction. Let ϕ be defined by (6.2) and, for any $n \geq 0$, let I_n be the result of an arbitrary quadrature formula (2.2) with nodes in $[-1, 1]$. Then the error is given by*

$$(6.5) \quad I - I_n = \frac{1}{2\pi i} \int_{\Gamma} f(z)(\phi(z) - r_n(z)) dz,$$

where r_n is defined by (6.4), and hence

$$(6.6) \quad |I - I_n| \leq \frac{1}{2\pi} |\Gamma| \|f\|_{\Gamma} \|\phi - r_n\|_{\Gamma},$$

where $|\Gamma|$ is the arc length of Γ and $\|\cdot\|_{\Gamma}$ is the supremum norm over Γ .

This theorem shows that quadrature over $[-1, 1]$ is associated with a problem in rational approximation: how closely can $\phi(z) = \log((z+1)/(z-1))$ be approximated in a region of the complex plane near $[-1, 1]$ by a rational function of type $(n, n+1)$, i.e., with numerator of degree $\leq n$ and denominator of degree $\leq n+1$? Note that $\phi(\infty) = 0$, so it makes sense that the degree of the denominator exceeds that of the numerator. Any sum (6.4) with nonzero weights w_k and distinct nodes x_k defines a rational function of type $(n, n+1)$ and, conversely, if a rational function of type $(n, n+1)$ is expanded in partial fractions, it will have the form (6.4) provided it has $n+1$ distinct poles.

The idea of relating quadrature formulas to rational functions goes back to Gauss himself [25, 26]. In fact, Gauss's original derivation is closer to this way of thinking than many of the subsequent developments framed in terms of orthogonal polynomials by Jacobi, Christoffel, Stieltjes, and others. Gauss quadrature formulas have a wonderfully simple characterization in terms of rational approximants.

THEOREM 6.2. *For any $n \geq 0$, the rational function r_n of (6.4) associated with the $(n+1)$ -point Gauss quadrature formula is the type $(n, n+1)$ Padé approximant to $\phi(z)$ at $z = \infty$. The function $\phi - r_n$ has a zero of order exactly $2n+3$ at $z = \infty$.*

Proof. Gauss derived this result by continued fractions [25]; see [26, section 1.2]. $\square$

A Padé approximant, as always, is defined as the unique rational function of the specified type that matches the Taylor expansion of the given function at the given point (here ∞) to as high an order as possible [3]. This terminology was not available in Gauss's day, but Theorem 6.2 is his nonetheless.

What about Clenshaw–Curtis quadrature?

We begin with an experimental comparison. Figure 4 shows errors $|(\phi - r_n)(z)|$ as contour plots in the complex plane for Newton–Cotes, Gauss, and Clenshaw–Curtis formulas with $n = 8$ and 16. In each case the innermost contour corresponds to $|\phi(z) - r_n(z)| = 1$ and, moving outward, the levels are $10^{-1}, 10^{-2}, \dots, 10^{-13}$. As $z \rightarrow \infty$, all the approximations become very good. Figures like these were drawn previously by Takahasi and Mori [54], for Gauss quadrature but not Clenshaw–Curtis, as part of an important series of papers on quadrature formulas and their enhancement by changes of variables like the double exponential transformation. The only other publication that I know of where figures like these appear is [56], for quadrature over $(-\infty, 0]$.

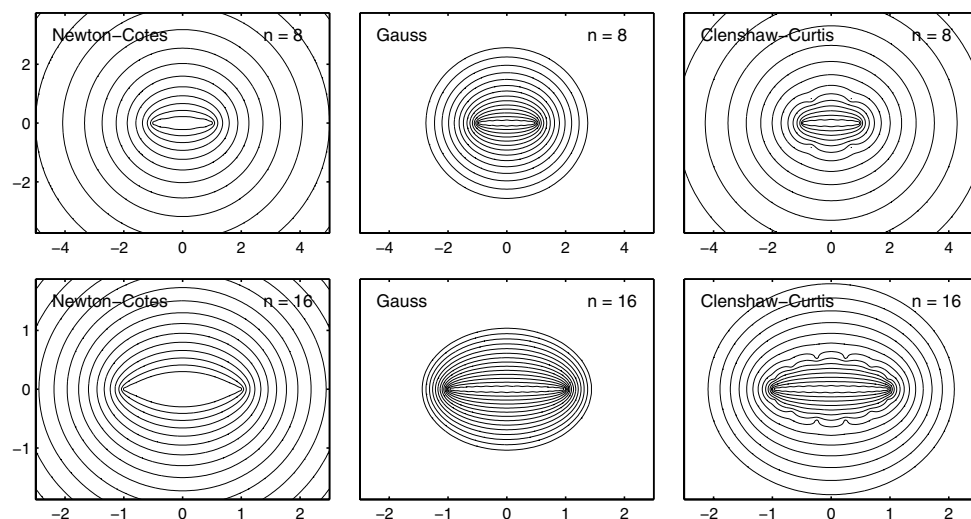


Fig. 4 Contour plots of the error $|\phi(z) - r_n(z)|$ for the rational approximations associated with Newton-Cotes, Gauss, and Clenshaw-Curtis quadrature formulas with $n = 8$ and 16 . The contours, from inner to outer, correspond to $10^0, 10^{-1}, \dots, 10^{-13}$. The Gauss and Clenshaw-Curtis errors are virtually identical near $[-1, 1]$, explaining their similar accuracy for functions that are not analytic in a large region of the plane.

What can we infer from such images? Suppose f is analytic in a big region of the complex plane and does not grow too fast as $|z|$ increases. Then (6.6) implies that we may get a small bound on $|I - I_n|$ by integrating over one of the contour levels far from $[-1, 1]$. In the extreme case, suppose f is a polynomial of degree k . Then Theorem 6.2 implies that for Gauss quadrature, $(\phi - r_n)f$ has a zero of order $2n + 3 - k$ at ∞ . If this number is at least 2, i.e., $k \leq 2n + 1$, then by taking circular contours Γ of radius $R \rightarrow \infty$ we get arbitrarily small bounds on $I - I_n$, confirming the exactness of Gauss quadrature for polynomials of degree $\leq 2n + 1$. By reversing the argument we can see that for any interpolatory formula, including Clenshaw-Curtis or Newton-Cotes, $\phi - r_n$ must have a zero of order at least $n + 2$ at ∞ . If it did not, taking a large circular contour Γ in (6.5) would imply a nonzero error $I - I_n$ for some monomial $f(x) = x^k$ with $k \leq n$.

At the other extreme, suppose f is analytic in only a small neighborhood of $[-1, 1]$ or not analytic at all. Then Figure 4 suggests that Newton-Cotes quadrature will not converge at all, which is indeed the case. But the remarkable thing is what it suggests about Clenshaw-Curtis quadrature. The figure reveals that near $[-1, 1]$, the Clenshaw-Curtis rational function is essentially indistinguishable from the Gauss rational function as an approximation to $\phi(z)$. Thus we can expect that for a function f of this kind, the two quadrature formulas should have about the same accuracy—exactly what was observed in the last three sections.

Figure 5 shows a close-up for the larger value $n = 64$. On the left we see contours for Gauss quadrature, and on the right, Clenshaw-Curtis. Now the pattern is even more striking. At $z = \infty$, r_n interpolates ϕ to order 131 for Gauss quadrature and only order 67 for Clenshaw-Curtis (since n is even, we get one more than the minimal order $n + 2$), and that is why the outer curves are twice as far apart on the right as on the left. The “missing” 64 interpolation points for Clenshaw-Curtis quadrature,

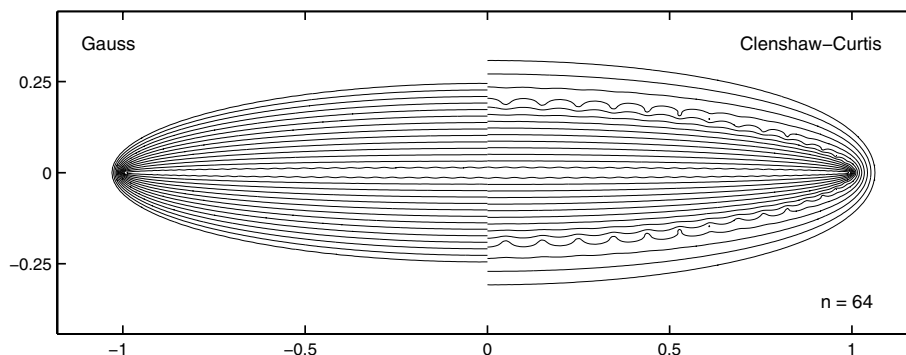


Fig. 5 Contour plot of the error $|\phi(z) - r_n(z)|$ as in Figure 4, but for $n = 64$ and with Gauss on the left half of the image, Clenshaw–Curtis on the right. The scallops in the curves on the right reveal the presence of $n - 2$ finite zeros of $\phi(z) - r(z)$ lying along a football-shaped curve surrounding $[-1, 1]$.

however, lie along a football-shaped curve at a certain distance from $[-1, 1]$ (actually there are 62 of them, not 64). Inside that curve, the two sides of the figure are almost indistinguishable.

The following is a first step towards explaining this effect.

THEOREM 6.3. *Let r_n be the rational function (6.4) associated with any interpolatory quadrature formula (2.2) with positive weights. Then the function $\phi(z) - r_n(z)$ has at least $n + 2$ zeros at $z = \infty$ and exactly $2n + 3$ zeros in $\mathbb{C} \cup \{\infty\} \setminus [-1, 1]$ if the points ± 1 are not quadrature nodes, or $2n + 1$ zeros if these points are nodes.*

Proof. A proof of the assertion about $z = \infty$ was sketched above. The result about total count of zeros in $\mathbb{C} \cup \{\infty\} \setminus [-1, 1]$ can be established by applying the argument principle to $\phi(z) - r(z)$ over a contour Γ that encloses $[-1, 1]$ sufficiently tightly. Each quadrature node in $(-1, 1)$ contributes a decrease in argument by 2π as Γ passes nearby in the upper half-plane, then another decrease by 2π as it passes nearby in the lower half-plane, and the final contributions come from $\phi(z)$, whose argument decreases by π near $x = 1$ and again by π near $x = -1$. We omit the details. $\square$

Thus we see that all interpolatory quadrature formulas with positive weights correspond to rational interpolants to $\phi(z)$ of order approximately $2n$. What distinguishes Gauss quadrature is that all the interpolation takes place at the single point ∞ . The rational functions associated with other quadrature formulas can be interpreted as examples of what approximation theorists call *multipoint Padé approximants* [3].

Figure 6 explores these interpolation points further, showing the $n - 2$ zeros of $\phi - r_n$ (computed numerically) for four values of n . (When n is odd the number of finite zeros increases to $n - 1$ and the number of zeros at $z = \infty$ shrinks from $n + 3$ to $n + 2$.) Two of the interpolation points are real, lying at approximately $\pm(1 + 1.7n^{-2})$, and the remainder lie in complex conjugate pairs along a curve whose height shrinks at the rate $O(n^{-1} \log n)$. Some of the details of this behavior are analyzed in a companion paper with Weideman [60]. In particular it would appear that a theorem along the following lines is likely to be valid, which would make an interesting companion to Theorem 5.1: the bound (4.14) applies to Clenshaw–Curtis as well as Gauss quadrature so long as $n \leq n_\rho$, where n_ρ is a number that approaches ∞ as ρ shrinks to 1.

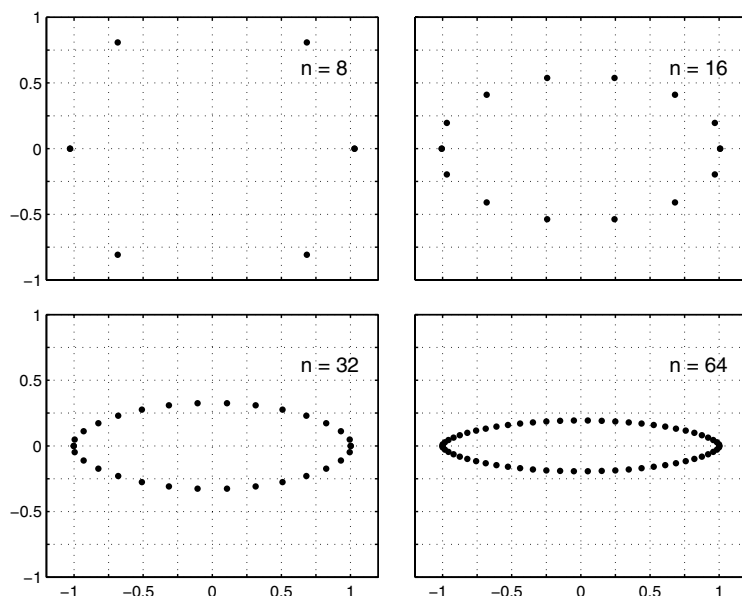


Fig. 6 Points in the complex plane at which $\phi(z) - r_n(z) = 0$ for Clenshaw–Curtis quadrature with $n = 8, 16, 32, 64$. In addition there are $n + 3$ interpolation points at $z = \infty$.

We shall close with an illustration of the implications of Figure 6. Consider again the fourth function of Figure 2, $f(x) = 1/(1 + 16x^2)$. This function is analytic on $[-1, 1]$ but has poles at $\pm i/4$. In Figure 6 we see that the curve of interpolation points crosses $\pm i/4$ somewhere between $n = 32$ and $n = 64$. This suggests that for values of n larger than about 50, the rate of convergence of the Clenshaw–Curtis formula will be cut in half. To test this prediction, Figure 7 extends the earlier computation to $n = 120$ instead of $n = 30$. As expected, there is a kink in the convergence curve at $n \approx 50$. Based on an early draft of this paper, an exact analysis of this example has been developed by Weideman and Trefethen [60] and related estimates for more general functions have been provided by Elliott, Johnston, and Johnston [17]. A previous analysis of convergence for functions of this form caught the asymptotic behavior, where Gauss beats Clenshaw–Curtis, but not the transient and the kink [6, Theorems 81 and 89].

7. Conclusion. Gauss quadrature, one of the jewels of numerical analysis, is a beautiful and powerful idea. Yet the Clenshaw–Curtis formula has essentially the same performance for most integrands and can be implemented effortlessly by the FFT. Figures 4 and 5 offer a visual explanation of why these two quadrature formulas are so close. Theorems 4.5 and 5.1 make the comparison quantitative.

The observations of this paper may have implications for spectral methods for the numerical solution of ordinary and partial differential equations [7], where the debate between Chebyshev ($\approx$ Clenshaw–Curtis) and Legendre ($\approx$ Gauss) methods is perennial.

Along the way to Theorem 5.1 on Clenshaw–Curtis quadrature, we have derived some possibly new results from the consideration of Chebyshev expansions of smooth functions: on decay of Chebyshev coefficients (see (4.6)), on near-optimality of truncated Chebyshev series (see (4.10) and (4.12)), on Jackson-type bounds for poly-

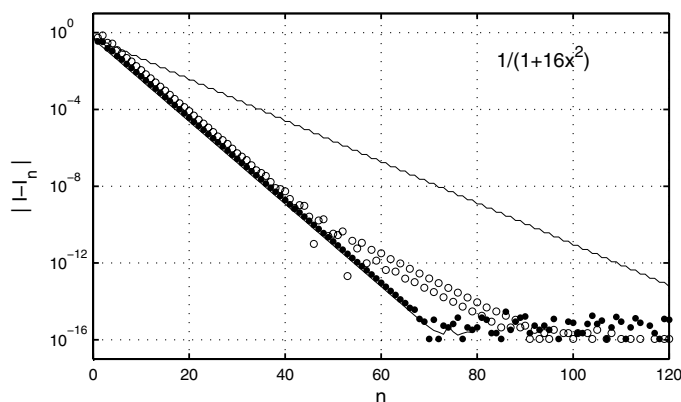


Fig. 7 Repetition of the fourth panel of Figure 2, but now to the higher value $n = 120$. Near $n = 50$, the rate of convergence is cut abruptly in half as the curve of interpolation points of Figure 6 crosses the poles of f at $\pm i/4$.

nomial best approximations (footnote 6), and on the accuracy of Gauss quadrature (see (4.13)). We do not claim, however, that any of these results are the best possible, and, indeed, as pointed out in section 4, they can probably be sharpened with the use of techniques of [13] and [14].

Though some results are given in [17] and [60], we hope that others will take up the challenge of explaining Figures 5–7 more fully, perhaps by making use of the potential theoretic notion of “balayage” [24]. The idea here would be that the differences $\phi - r_n$ for Gauss and Clenshaw–Curtis quadrature are controlled by the zeros of this function, analogous to point charges. For Gauss quadrature all the charges lie at $z = \infty$, while for Clenshaw–Curtis half of them have been “swept out” from ∞ to a finite curve; but the potentials generated in both cases are essentially the same near $[-1, 1]$.

This article has questioned the factor of 2 in efficiency that is widely thought to separate Clenshaw–Curtis and Gauss quadrature. Quite another and independent matter is the factor of $\pi/2$ that both methods lose due to clustering of quadrature nodes near endpoints, in comparison with the trapezoid rule for a periodic integrand. As mentioned in footnote 4, new “transplanted Gauss” and “transplanted Clenshaw–Curtis” methods based on conformal maps are proposed in [33] for recovering this factor of $\pi/2$ (joint work with Nicholas Hale). Combining the results of that paper and this one, we find that transplanted Clenshaw–Curtis quadrature is both faster than Gauss quadrature to implement and more accurate.

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